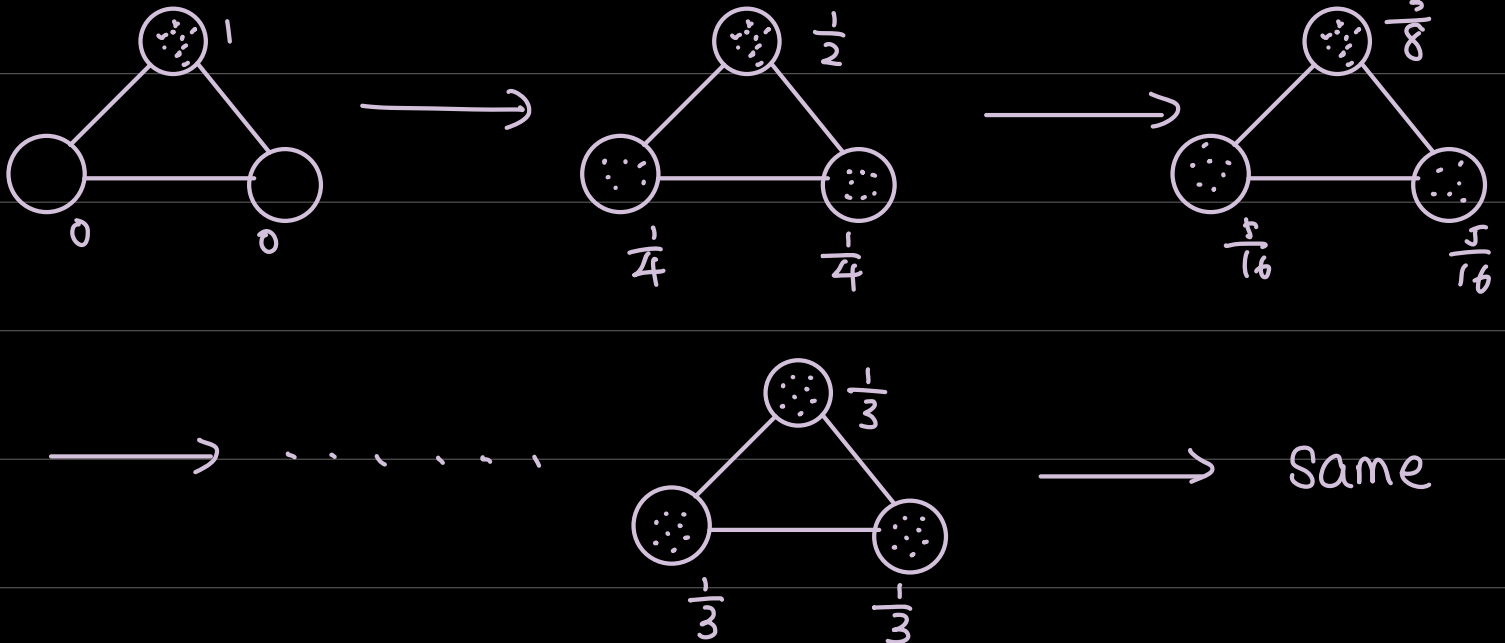


# Markov Chain

random walk in state space

1 million people



Population migration

probability  $\frac{1}{2}$ , stay, probability  $\frac{1}{4}$  to one other state.

Arrow of time, coverage to stationary distribution

Entropy (randomness) increases

Notation :

$X_t$  : State at time  $t$

$p^{(t)}$  : distribution of  $X_t$  (population distribution)

$p^{(t)}$	$p_1^{(t)}$	$p_2^{(t)}$	$p_3^{(t)}$
$t=0$	1	0	0
$t=1$	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
$t=2$	$\frac{3}{8}$	$\frac{5}{16}$	$\frac{5}{16}$
$\vdots$	$\vdots$	$\vdots$	$\vdots$
$t \rightarrow \infty \pi$	$\frac{1}{3}\pi_1$	$\frac{1}{3}\pi_2$	$\frac{1}{3}\pi_3$

$$p^{(t)} \xrightarrow[t \rightarrow \infty]{} \pi$$

"arrow of time"

Key / driving force :

Transition Probability

$$K_{ij} = P(X_{t+1} = j \mid X_t = i)$$

$$\text{or } [K(x, y) = P(X_{i+1} = y \mid X_i = x)]$$

transition matrix:

$j$	1	2	3
$i$			
1	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
2	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$
3	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$

# Master Equation

$$\begin{aligned}
 P_j^{(t+1)} &= P(X_{t+1} = j) \quad \rightarrow \text{\# of ppl in } j \text{ at } t+1 \\
 &= \sum_i P(X_{t+1} = j \ \& \ X_t = i) \quad \rightarrow \text{\# in } j \text{ at } t+1 \text{ and in } i \text{ at } t \\
 &= \sum_i P(X_t = i) P(X_{t+1} = j | X_t = i) \quad \rightarrow \text{\# in } i \text{ at } t \\
 &= \sum_i P_i^{(t)} K_{ij} \quad \begin{array}{l} \downarrow \\ \text{fraction of these in } i \\ \text{who will go to } j \end{array}
 \end{aligned}$$

So the master equation is:

$$P_j^{(t+1)} = \sum_i P_i^{(t)} K_{ij}$$

eigen-analysis

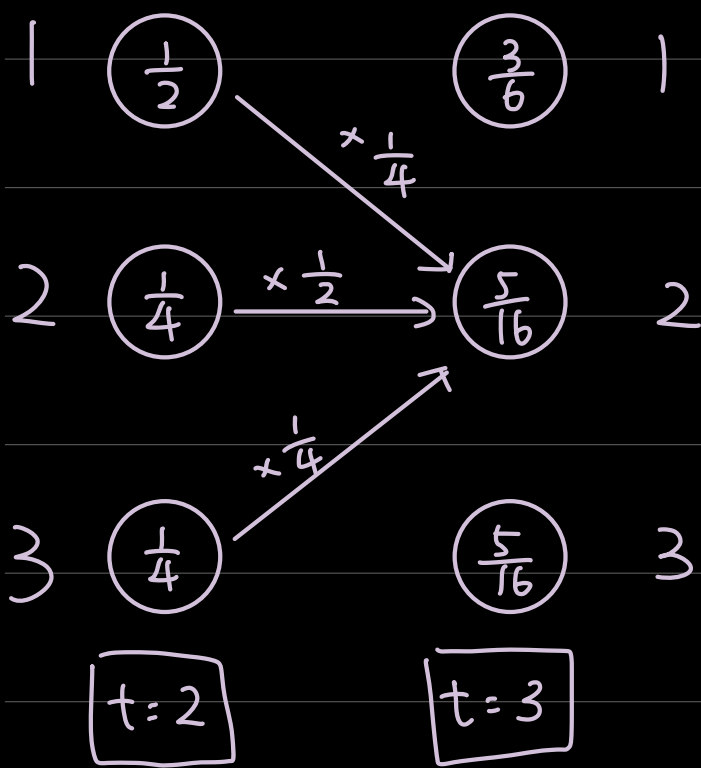
$$P^{(t)} \xrightarrow{K} P^{(t+1)} \quad \text{converge} \nearrow \pi$$

$$P^{(t+1)} = P^{(t)} K \Rightarrow P^{(t)} = P^{(0)} K^t$$

$$\begin{bmatrix} \frac{3}{8} & \frac{5}{16} & \frac{5}{16} \end{bmatrix} = \begin{bmatrix} \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \end{bmatrix}$$

$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$
$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$

Interpret the master equation



$P_i^{(t)}$  = # of people in  $i$  at  $t$

$K_{ij}$  = fraction of people in  $i$  who will go to  $j$

$P_i^{(t)} K_{ij}$  = # of people in  $i$  at  $t$  who will go to  $j$

$\sum P_i^{(t)} K_{ij}$  = # of people who will end up in  $j$  at  $t+1$

$\approx$  equilibrium

## Stationary $\pi$

$$\pi = \pi K$$

$$\pi_j = \sum_i \pi_i K_{ij} \rightarrow \text{Overall balance}$$

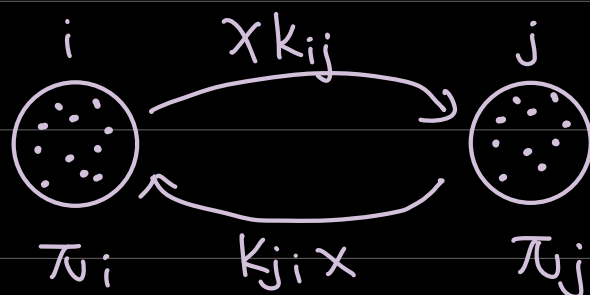
②

Detailed balance  $\forall i, j$

① condition  
stronger condition:

satisfies overall balance when it's satisfied.

$$\pi_i K_{ij} = \pi_j K_{ji}$$



Page Rank : know  $K$ , want  $\pi$

Start from  $p^{(0)}$  (e.g. uniform)

$$p^{(0)} \xrightarrow{K} p^{(1)} \xrightarrow{K} p^{(2)} \xrightarrow{K} \dots \xrightarrow{K} p^{(T)} \approx \pi$$

$K$  is a stochastic matrix, the sum of each row = 1

Why recursive : more stable, accurate measure of popularity

MCMC : know  $\pi$  <sup>(target distribution)</sup>, want  $K$  <sup>(iterative algorithm)</sup>  
so that  $\pi = \pi K$  <sup>(design principle)</sup> <sup>to implement</sup>

Start from  $X_0 \sim p^{(i)}$

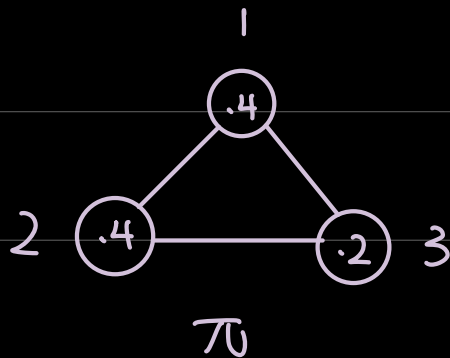
$$X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow \dots \rightarrow X_t \rightarrow X_{t+1} \rightarrow \dots \rightarrow X_T$$

Approx:  $X_T \sim \pi$

# Metropolis algorithm

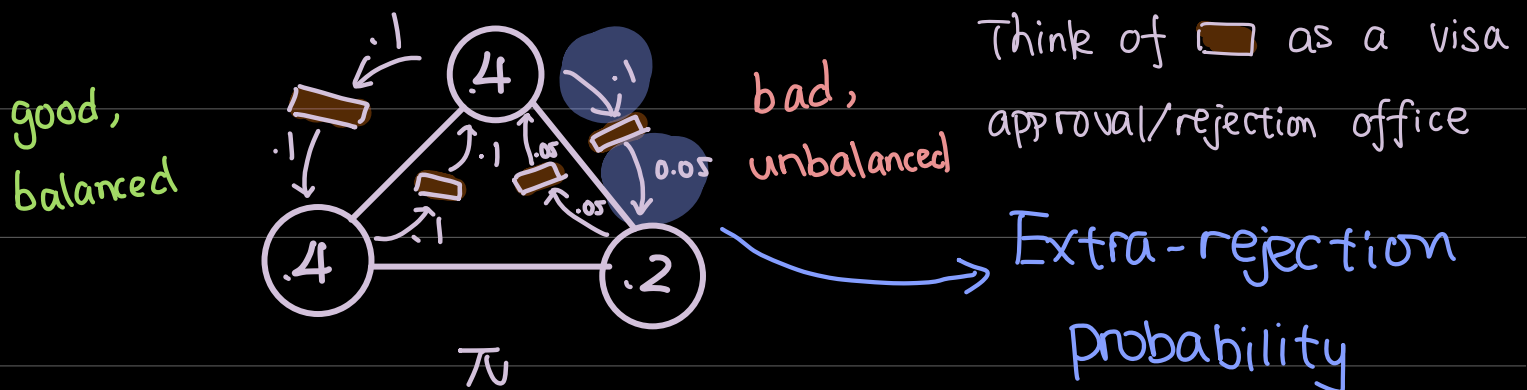
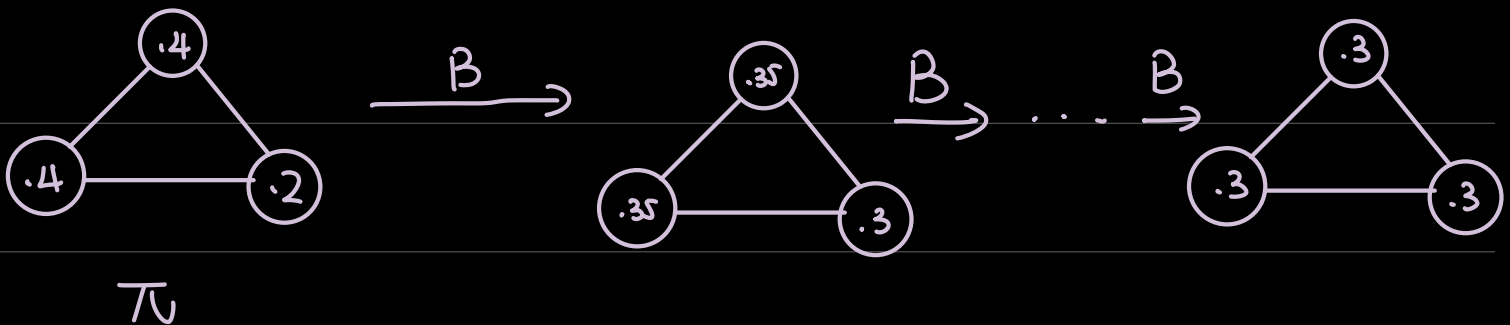
1950s Metropolis, Roseblatt<sup>2</sup>, Teller<sup>2</sup>

Basic idea

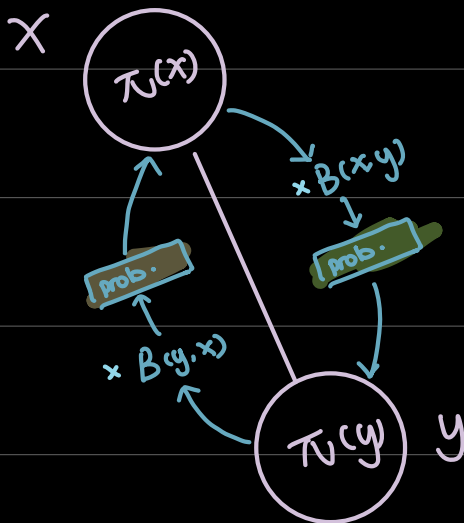


Base Chain:  $B$

$i \backslash j$	1	2	3
1	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
2	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$
3	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$



General:



Algorithm:

(1) If  $\pi(x)B(x,y) > \pi(y)B(y,x)$

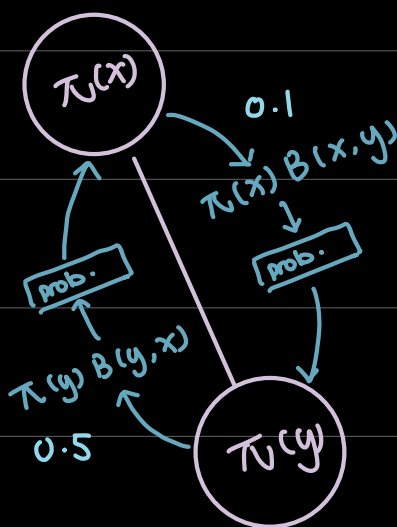
$$\text{prob}(\text{accept } x \rightarrow y) = \frac{\pi(y)B(y,x)}{\pi(x)B(x,y)}$$

(2) If  $\pi(x)B(x,y) \leq \pi(y)B(y,x)$

$$\text{prob}(\text{accept } x \rightarrow y) = 1$$

(3) Combine (1), (2):

$$\text{prob}(\text{accept } x \rightarrow y) = \min\left(1, \frac{\pi(y)B(y,x)}{\pi(x)B(x,y)}\right)$$



$$X_t = x$$

$$X_{\text{propose}} \sim B(x, y) = P(X_{\text{propose}} = y | X_t = x)$$

$$\text{let } X_{t+1} = \begin{cases} X_{\text{propose}} & \text{with } p = \min\left(1, \frac{\pi(y) B(y, x)}{\pi(x) B(x, y)}\right) \\ X_t = x & \text{with } 1-p \end{cases}$$

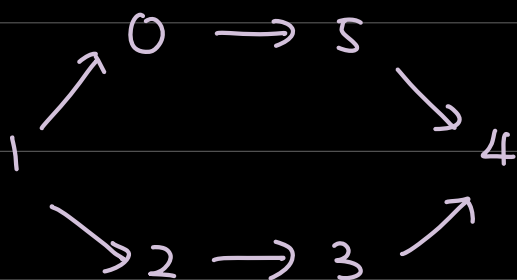
Special Case :

Metropolis algorithm :  $B(x, y) = B(y, x)$

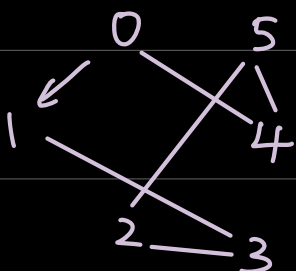
$$p = \min\left(1, \frac{\pi(y)}{\pi(x)}\right)$$

Hastings Extension : general  $B$

# Traveling Salesman Problem



$X$ : path  $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 0$



another  $X$  :  $0 \rightarrow 1 \rightarrow 3 \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 0$



$\epsilon(x)$  = total length of path  $x$ .

$$\pi(x) = \frac{1}{Z} e^{-\epsilon(x)}$$

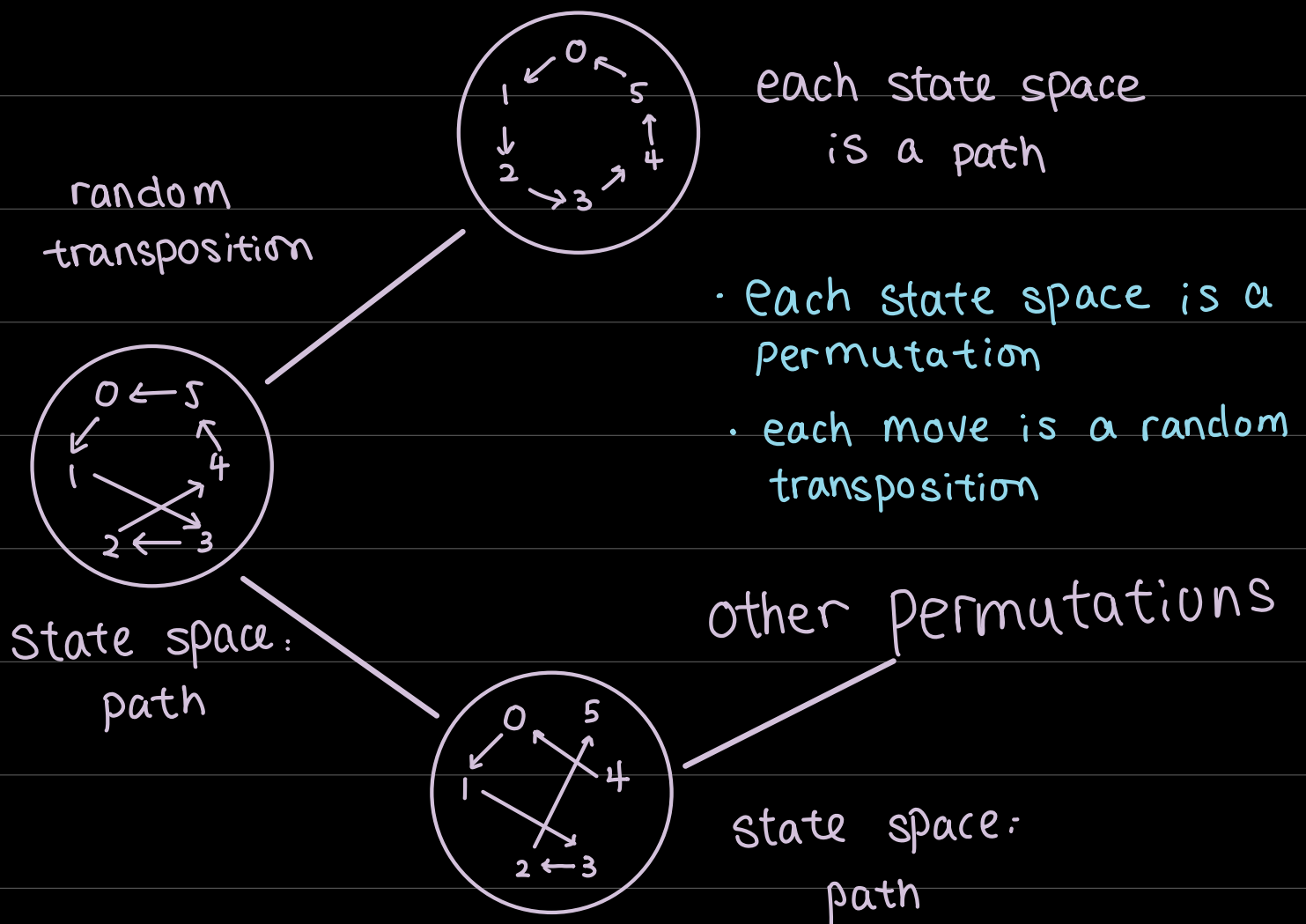
$$Z = \sum_x e^{-\epsilon(x)}$$

Gibbs distribution / Boltzman distribution

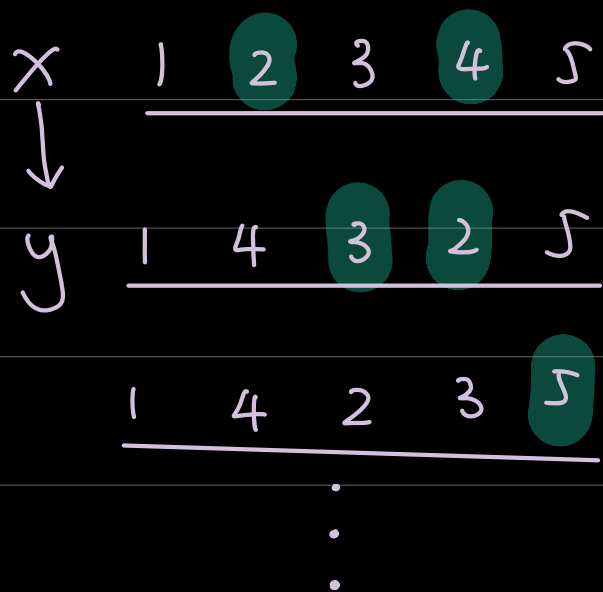
Searching thru such huge space is extremely difficult

sampling and optimization are related here

State space (n factorial)



## Base Chain



Randomly pick 2 cities  
& exchange their positions,

Propose change by random  
transposition.

accept change with

$$p = \min \left( 1, \frac{\pi(y)}{\pi(x)} \right)$$

$$= \min \left( 1, e^{\frac{1}{2}(\epsilon(x) - \epsilon(y))} \right)$$

↑  
here  $\frac{1}{2}$  gets cancelled

recall

$$\pi(-) = \frac{1}{2} e^{-\epsilon(-)}$$

$$X_0 \longrightarrow X_1 \longrightarrow X_2 \longrightarrow X_3 \longrightarrow \dots$$

(1 2 3 4 5)      (1 4 3 2 5)      (1 4 2 3 5)      (1 4 2 3 5)

→  $X_{1000}$  we end up w/ a good permutation

$$\text{Approx: } X_{1000} \sim \pi(x)$$